

# Counting modulo infinite monoids and $\Delta_0$ -definability

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## Plan

- What is  $\Delta_0$ -definability and why counting modulo (finite) monoids?

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- A new result about  $SL(2, \mathbb{N})$

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- A new result about  $SL(2, \mathbb{N})$
- Counting modulo infinite monoids
- A new result about  $SL(2, \mathbb{Z})$
- Conclusion and further work

What is  $\Delta_0$ -definability?

$x$  is not prime nor 0 nor 1

$$(\exists u)_{u < x} (\exists v)_{v < x} (x = uv)$$

## What is $\Delta_0$ -definability?

$x$  is not prime nor 0 nor 1

$$(\exists u)_{u < x} (\exists v)_{v < x} (x = uv)$$

$\Delta_0$ -definability is essentially definability with a formula in the language of arithmetic where the quantified variables are bounded by terms



What is  $\Delta_0$ -definability?

Major open problem

Find a "simple" arithmetical relation

NOT  $\Delta_0$ -definable

Here "simple" is a non defined meta-mathematical notion!

What is  $\Delta_0$ -definability?

Open example

The relation *y is the n-th prime number*

IS NOT KNOWN TO BE  $\Delta_0$ -definable

What is  $\Delta_0$ -definability?

Open example

The relation *y is a prime number of even index*

IS NOT KNOWN TO BE  $\Delta_0$ -definable

## Why counting modulo (finite) monoids?

The relation  $y$  is a *prime number of even index*

is defined by  $(y \text{ is prime}) \wedge (f(y) = 0)$  where

$f$  is recursively defined by

$$\left\{ \begin{array}{l} f(0) = 0 \\ f(i+1) = f(i) + 1 \pmod{2} \text{ if } i \text{ is prime} \\ f(i+1) = f(i) \text{ if } i \text{ is not prime} \end{array} \right.$$

## Why counting modulo (finite) monoids?

The relation  $y$  is a prime number of even index

is defined by

$$(y \text{ is prime}) \wedge (g(0) +_2 g(1) +_2 \dots +_2 g(y) = 0)$$

$$\text{where } \begin{cases} g(i) = 1 \in \mathbb{Z}/2\mathbb{Z} \text{ if } i \text{ is prime} \\ g(i) = 0 \in \mathbb{Z}/2\mathbb{Z} \text{ if } i \text{ is not prime} \end{cases}$$

$+_2$  is the sum modulo 2.

## Why counting modulo (finite) monoids?

$\Delta_0^{\sharp G}$  is defined by adding to the definition of  $\Delta_0$   
the following closure under counting modulo  $G$   
where  $G$  is a finite monoid:

## Why counting modulo (finite) monoids?

Suppose that  $\begin{cases} \mathbb{N} & \rightarrow & G \\ i & \rightarrow & g(i) \end{cases}$  is s.t.

for all  $g \in G$ , the relation  $g(i) = g$  is  $\Delta_0^{\#G}$ -definable

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THEN

for all  $g \in G$ , the relation  $g(0) +_G g(1) +_G \dots +_G g(y) = g$   
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Counting modulo  $\mathbb{N}$

Strong counting modulo  $\mathbb{N}$

## Counting modulo $\mathbb{N}$

## Strong counting modulo $\mathbb{N}$

Suppose that  $\begin{cases} \mathbb{N} & \rightarrow & \mathbb{N} \\ i & \rightarrow & g(i) \end{cases}$  is s.t.  
the relation  $z = g(i)$  is  $\Delta_0^{\#\mathbb{N}}$ -definable

## Counting modulo $\mathbb{N}$

## Strong counting modulo $\mathbb{N}$

Suppose that  $\begin{cases} \mathbb{N} & \rightarrow & \mathbb{N} \\ i & \rightarrow & g(i) \end{cases}$  is s.t.

the relation  $z = g(i)$  is  $\Delta_0^{\#\mathbb{N}}$ -definable

THEN

the relation  $g(0) + g(1) + \dots + g(y) = z$   
is  $\Delta_0^{\#\mathbb{N}}$ -definable

Counting modulo  $\mathbb{N}$

Weak counting modulo  $\mathbb{N}$

## Counting modulo $\mathbb{N}$

## Weak counting modulo $\mathbb{N}$

Suppose that  $\begin{cases} \mathbb{N} & \rightarrow & \{0; 1\} \\ i & \rightarrow & g(i) \end{cases}$  is s.t.  
the relation  $1 = g(i)$  is  $\Delta_0^{\#\mathbb{N}}$ -definable

## Counting modulo $\mathbb{N}$

## Weak counting modulo $\mathbb{N}$

Suppose that  $\begin{cases} \mathbb{N} & \rightarrow & \{0; 1\} \\ i & \rightarrow & g(i) \end{cases}$  is s.t.

the relation  $1 = g(i)$  is  $\Delta_0^{\#\mathbb{N}}$ -definable

THEN

the relation  $g(0) + g(1) + \dots + g(y) = z$   
is  $\Delta_0^{\#\mathbb{N}}$ -definable

Counting

Known facts

Counting

Known facts

JAF's folklaw:

Previous definitions of counting modulo  $\mathbb{N}$  are equivalent



## Counting

## Known facts

JAF's folklaw+ **Theorem** (Clote 95)

$$\Delta_0^{\#_{\mathbb{Z}/2\mathbb{Z}}} \subseteq \Delta_0^{\#_{\mathbb{Z}}} \subseteq \Delta_0^{\#_{\sigma_5}}$$

$\sigma_5$  is the permutation group over five elements.

$$\begin{array}{l} \Delta_0^{\#_{\mathbb{Z}/2\mathbb{Z}}} \\ \Delta_0^{\#_{\mathbb{Z}/3\mathbb{Z}}} \end{array} \subseteq \Delta_0^{\#_{\mathbb{Z}/6\mathbb{Z}}}$$

## A result about $SL(2, \mathbb{N})$

### Theorem

$$\text{Let } \begin{cases} \mathbb{N} \rightarrow & SL(2, \mathbb{N}) \\ i \rightarrow & \begin{pmatrix} a(i) & b(i) \\ c(i) & d(i) \end{pmatrix} \end{cases}$$

where  $\begin{pmatrix} u & v \\ w & z \end{pmatrix} = \begin{pmatrix} a(i) & b(i) \\ c(i) & d(i) \end{pmatrix}$  is  $\Delta_0$ -definable

## A result about $SL(2, \mathbb{N})$

### Theorem

Then

$$\begin{pmatrix} u & v \\ w & z \end{pmatrix} = \begin{pmatrix} a(0) & b(0) \\ c(0) & d(0) \end{pmatrix} \begin{pmatrix} a(1) & b(1) \\ c(1) & d(1) \end{pmatrix} \cdots \begin{pmatrix} a(y) & b(y) \\ c(y) & d(y) \end{pmatrix}$$

is  $\Delta_0^{\sharp\mathbb{N}}$ -definable

A result about  $SL(2, \mathbb{N})$

Idea of proof

## A result about $SL(2, \mathbb{N})$

### Idea of proof

#### Lemma 1

$$\text{Let } M = \begin{pmatrix} a & b \\ c & d \end{pmatrix} \in SL(2, \mathbb{N})$$

There exist  $k(a, b, c, d)$  and  $\alpha(i, a, b, c, d)$  s.t.

$$M = \begin{pmatrix} 1 & \alpha(0, a, b, c, d) \\ 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ \alpha(1, a, \dots) & 1 \end{pmatrix} \cdots \begin{pmatrix} 1 & \alpha(2k(a, \dots), a, \dots) \\ 0 & 1 \end{pmatrix}$$

and  $\alpha(0) \geq 0$ ,  $\alpha(1) \geq 1$ , ...  $\alpha(2k-1) \geq 1$ ,  $\alpha(2k) \geq 0$ ,

and  $\alpha$  and  $k$  have  $\Delta_0$ -definable graph.

## A result about $SL(2, \mathbb{N})$

### Idea of proof

#### Lemma 2

Let  $\begin{pmatrix} a & b \\ c & d \end{pmatrix} \in SL(2, \mathbb{N})$ , and  $\alpha$  a function with a  $\Delta_0$ -definable graph  
and  $\alpha(0) \geq 0, \alpha(1) \geq 1, \dots, \alpha(2k-1) \geq 1, \alpha(2k) \geq 0,$

Then  $\begin{pmatrix} a & b \\ c & d \end{pmatrix} = \begin{pmatrix} 1 & \alpha(0) \\ 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ \alpha(1) & 1 \end{pmatrix} \cdots \begin{pmatrix} 1 & \alpha(2k) \\ 0 & 1 \end{pmatrix}$   
is  $\Delta_0$ -definable

## A result about $SL(2, \mathbb{N})$

### Idea of proof

#### Lemma 2bis

Let  $\begin{pmatrix} a & b \\ c & d \end{pmatrix} \in SL(2, \mathbb{N})$ , and  $\alpha$  a function with a  $\Delta_0$ -definable graph

and  $\alpha(0) \geq 0$ ,  $\alpha(1) \geq 0$ , ...  $\alpha(2k-1) \geq 0$ ,  $\alpha(2k) \geq 0$ ,

$$\text{Then } \begin{pmatrix} a & b \\ c & d \end{pmatrix} = \begin{pmatrix} 1 & \alpha(0) \\ 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ \alpha(1) & 1 \end{pmatrix} \cdots \begin{pmatrix} 1 & \alpha(2k) \\ 0 & 1 \end{pmatrix}$$

is  $\Delta_0^{\# \mathbb{N}}$ -definable

What could be counting modulo infinite monoids?

What could be counting modulo infinite matrix monoids?



What could be counting modulo infinite monoids?

What could be counting modulo infinite matrix monoids?

Let  $M(i) = (a_{u,v}(i))_{1 \leq u,v \leq d} \in G$

Suppose  $(z_{u,v})_{1 \leq u,v \leq d} = (a_{u,v}(i))_{1 \leq u,v \leq d}$  is  $\Delta_0^{\#G}$ -definable

What could be counting modulo infinite monoids?

What could be counting modulo infinite matrix monoids?

Let  $M(i) = (a_{u,v}(i))_{1 \leq u,v \leq d} \in G$

Suppose  $(z_{u,v})_{1 \leq u,v \leq d} = (a_{u,v}(i))_{1 \leq u,v \leq d}$  is  $\Delta_0^{\#G}$ -definable

THEN

$(z_{u,v})_{1 \leq u,v \leq d} = M(0)M(1)\dots M(y)$  is  $\Delta_0^{\#G}$ -definable

What could be counting modulo infinite monoïds?

What could be counting modulo infinite finitely generated monoïds?

Let  $G$  be a monoïd with a finite set of generators  $\Gamma = \{\gamma_1, \dots, \gamma_d\}$

What could be counting modulo infinite monoïds?

What could be counting modulo infinite finitely generated monoïds?

Let  $G$  be a monoïd with a finite set of generators  $\Gamma = \{\gamma_1, \dots, \gamma_d\}$

For all  $\begin{cases} N & \rightarrow & \Gamma \\ i & \rightarrow & g(i) \end{cases}$  s.t.

for all  $\gamma \in \Gamma$ , the relation  $g(i) = \gamma$  is  $\Delta_0^{\sharp G}$ -definable

THEN

for all  $\gamma \in \Gamma$ , the relation  $g(0).g(1)...g(z) = \gamma$  is  $\Delta_0^{\sharp G}$ -definable

What could be counting modulo infinite monoïds?

What could be counting modulo infinite finitely generated monoïds of matrices?

The first definition is stronger than the second

A result about  $SL(2, \mathbb{Z})$

Definition

What is  $\Delta_0$ -definability in  $\mathbb{Z}$ ?

A result about  $SL(2, \mathbb{Z})$

Definition

What is  $\Delta_0$ -definability in  $\mathbb{Z}$ ?

just consider a relative integer as couple in  $\{0; 1\} \times \mathbb{N}$

## A result about $SL(2, \mathbb{Z})$

### Definition

Hence,  $a : \mathbb{N} \mapsto \mathbb{Z}$  has a  $\Delta_0$ -definable graph

iff  $|a|$  has a  $\Delta_0$ -definable graph  
and  $sign(a(i)) = 1$  is a  $\Delta_0$ -definable relation



A result about  $SL(2, \mathbb{Z})$

Theorem

## A result about $SL(2, \mathbb{Z})$

### Theorem

$$\text{Let } \begin{cases} \mathbb{N} \rightarrow & SL(2, \mathbb{Z}) \\ i \rightarrow & \begin{pmatrix} a(i) & b(i) \\ c(i) & d(i) \end{pmatrix} \end{cases}$$

where  $a, b, c, d$  are polynomially bounded,

and  $\begin{pmatrix} u & v \\ w & z \end{pmatrix} = \begin{pmatrix} a(i) & b(i) \\ c(i) & d(i) \end{pmatrix}$  is a  $\Delta_0$ -definable relation

A result about  $SL(2, \mathbb{Z})$

Theorem

THEN

$$\begin{pmatrix} u & v \\ w & z \end{pmatrix} = \begin{pmatrix} a(0) & b(0) \\ c(0) & d(0) \end{pmatrix} \begin{pmatrix} a(1) & b(1) \\ c(1) & d(1) \end{pmatrix} \cdots \begin{pmatrix} a(y) & b(y) \\ c(y) & d(y) \end{pmatrix}$$

is  $\Delta_0^{\aleph_N}$  definable

A result about  $SL(2, \mathbb{Z})$

Idea of proof

Generalize the results about  $\Delta_0$ -definability concerning  
the standard euclidean algorithm  
to the least absolute remainder euclidean algorithm:

A result about  $SL(2, \mathbb{Z})$

Idea of proof

$$r_n = \beta_n r_{n+1} + r_{n+2}$$

where  $\frac{|r_{n+1}|}{2} < |r_{n+2}| \leq \frac{|r_{n+1}|}{2}$

## Conclusion and further work

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The proof does not extend to

$$SL(3, \mathbb{Z})$$

## Conclusion and further work

**Question 1:** Find a convenient definition for

### **Multiple Continued Fractions**

for solving the  $SL(3, \mathbb{Z})$  case.



## Conclusion and further work

**Question 2:** Compare both definitions of counting  
modulo infinite monoids  
in the case of finitely generated  
monoids of matrices

## Conclusion and further work

**Work:** Consider counting modulo infinite monoids  
in the case of a finite presentation  
with generators and relations